

Routed End-to-End Ethernet Network Proof of Concept

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May 21st, 2008

Agenda

- ✦ Introduction
- ✦ Ethernet and IEEE 802.1 (layer 2)
- ✦ IPv4 and IPv6 (layer 3)
- ✦ RE2EE
- ✦ Proof of Concept
- ✦ Results
- ✦ Conclusions

Introduction

- ✦ Problems in the contemporary Internet
 - ✦ IPv4 address space is running out of addresses
 - ✦ Routing tables are getting oversized
- ✦ IPv6 has been implemented, but deployment is not progressing
- ✦ Something needs to be done!

Ethernet

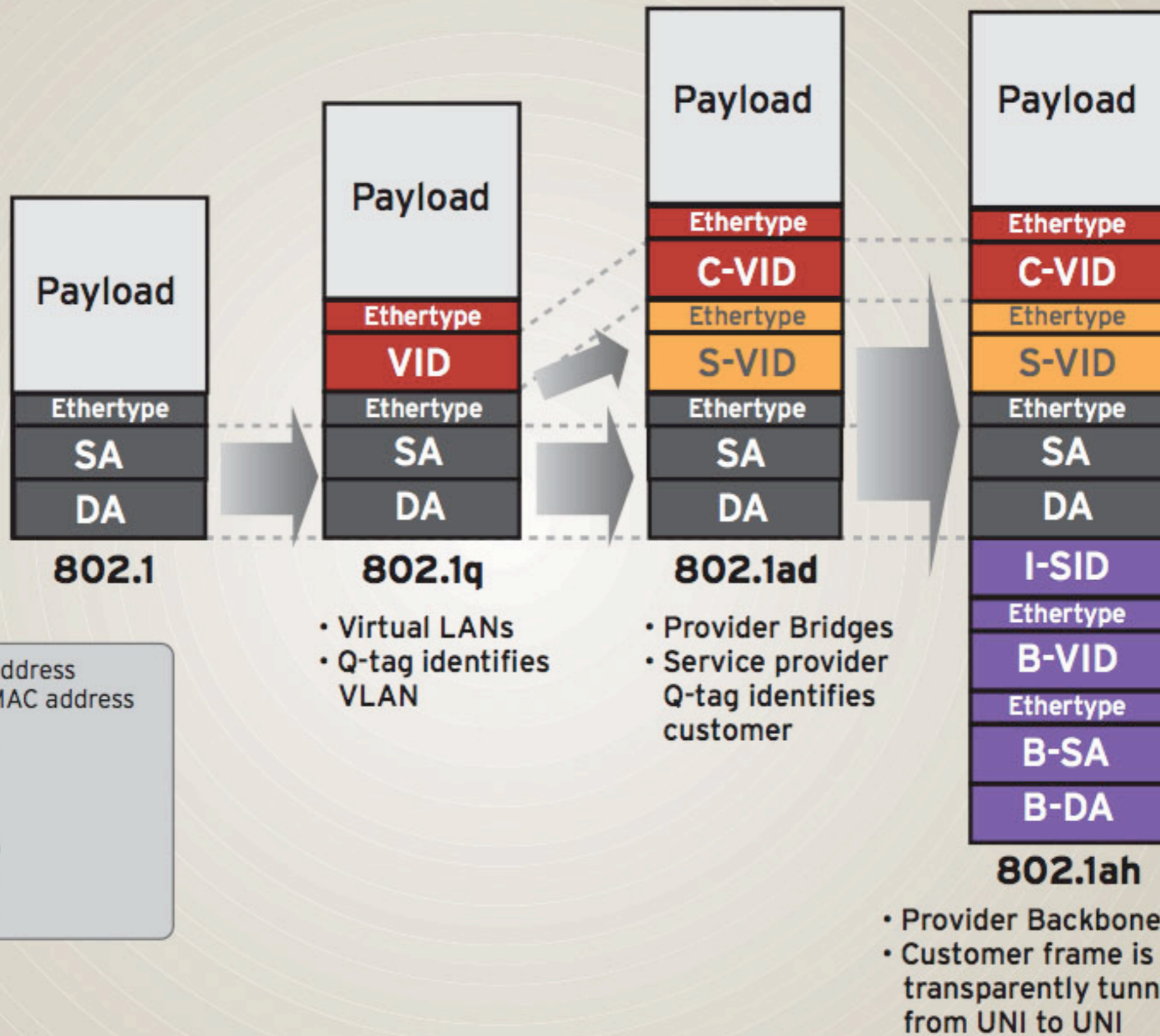
- ✦ Ethernet has a large address space with its MAC addresses
 - ✦ Not expected to be exhausted sooner than the year 2100
 - ✦ EUI-48 (MAC) could be even enlarged with EUI-64
- ✦ Ethernet is widely used and low-cost to implement
- ✦ Does not have hierarchy in addresses -> no efficient routing -> does not scale well

IEEE 802.1 (1/2)

- ✦ There are many standards and drafts to enhance the scalability of Ethernet networks
- ✦ 802.1Q - Virtual LANs
 - ✦ Q-tag adds 2nd tier to network's hierarchy
 - ✦ Inter-VLAN communication needed to be done using layer 3 routers
- ✦ 802.1ad - Provider Bridges
 - ✦ Q-in-Q added support for three-tiered hierarchy
 - ✦ Still supports only 4094 customer VLANs for ISPs

IEEE 802.1 (2/2)

- ✦ 802.1ah - Provider Backbone Bridges
 - ✦ MAC-in-MAC separates Ethernet network into customer and provider domains with complete isolation among their MAC addresses
 - ✦ It encapsulates the customer MAC header with a service provider MAC header
 - ✦ Supports a theoretical maximum of 16 million service instances
 - ✦ Still a draft and the newest version is 4.0



SA = Source MAC address
 DA = Destination MAC address
 VID = VLAN ID
 C-VID = Customer VID
 S-VID = Service VID
 I-SID = Service ID
 B-VID = Backbone VID
 B-DA = Backbone DA
 B-SA = Backbone SA

Development of Ethernet headers

IPv4 and IPv6

- ✦ IPv4 became the main protocol for Internet because of the hierarchical addresses, which allowed efficient routing
- ✦ Now IPv4 address space is running out of addresses
- ✦ Network Address Translation has been postponing the problem, but at the same time causing other problems
- ✦ IPv6 enlarges the address space
- ✦ IPv6 has been implemented, but deployment is not progressing

Routed End-to-End Ethernet

- ✦ The main idea is to use Ethernet instead of IP
- ✦ MAC addresses and NSAP addresses as network locators for hosts and servers
- ✦ Moving from layer 3 to layer 2 means that the technology must have routing capability
- ✦ Routing needs hierarchy in the addresses -> NSAP
- ✦ IS-IS could be the routing protocol for RE2EE

Network Structure in RE2EE

- ✦ Network Core

- ✦ RE2EE network can have many network cores

- ✦ Service Core

- ✦ Provides VPNs and Public Service Networks each in their own overlay
- ✦ Supports parallel networks such as Internet, packet TV network, P2P network

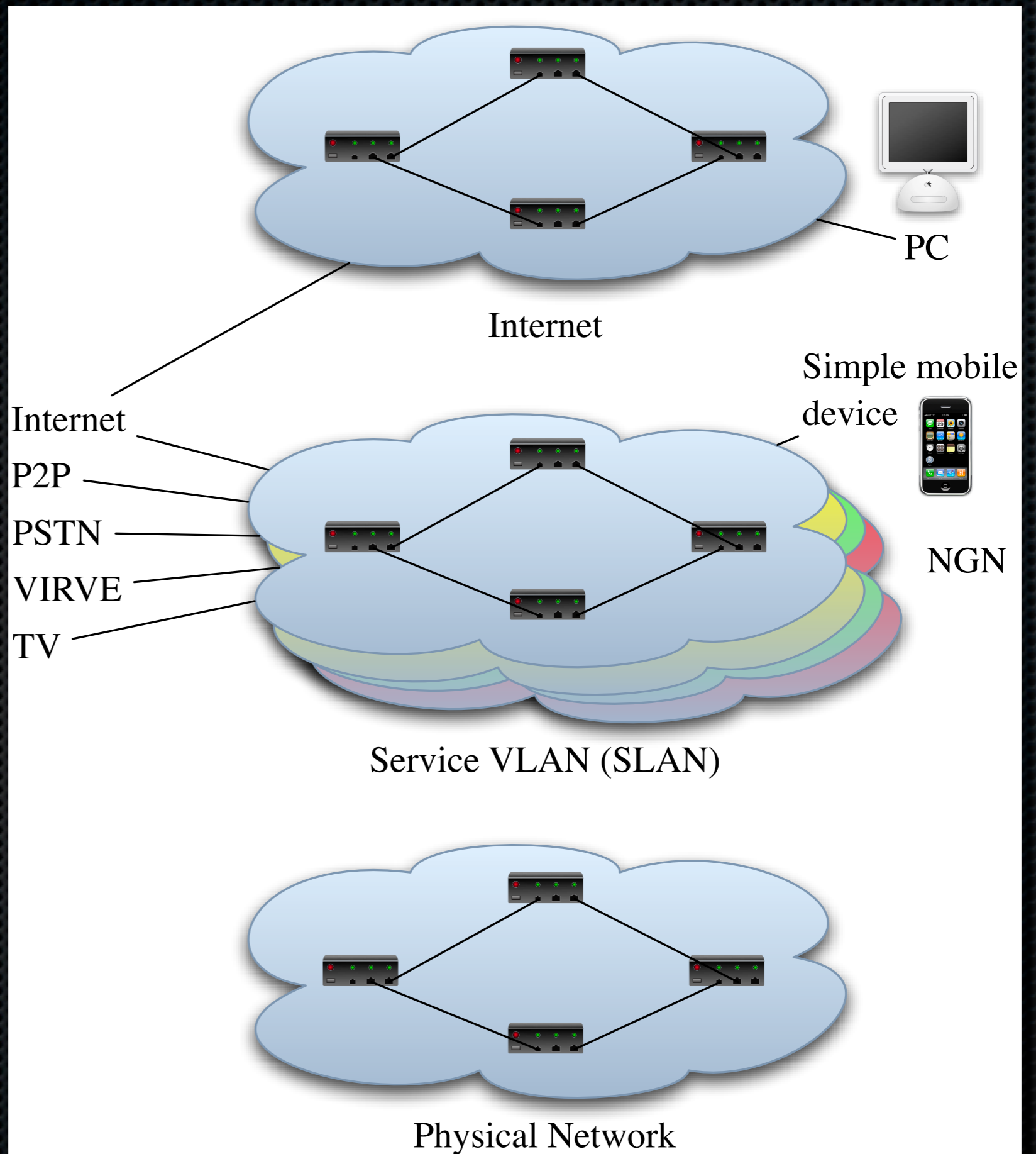
- ✦ Mobility Layer

- ✦ User Identity Layer

- ✦ Network Service Access Point (NSAP) addresses are used in the Provider side and MAC addresses are used in the user side

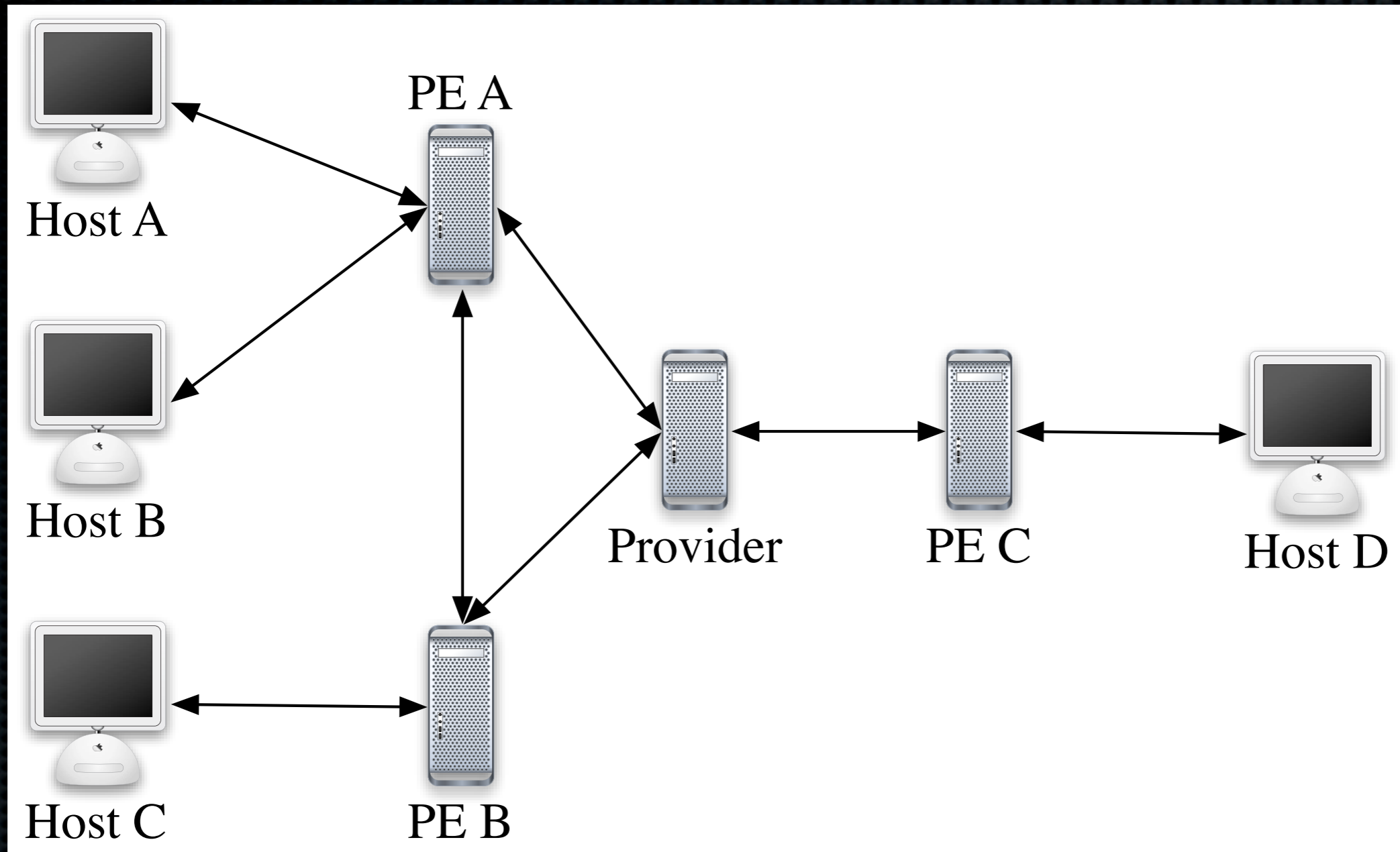
Service VLANs

- ✦ Internet
- ✦ P2P
- ✦ PSTN
- ✦ TV



Proof of Concept

- ✦ Proof of Concept does not implement all the features from the RE2EE concept
- ✦ Network is built on two PCs running Debian GNU/Linux
- ✦ Programming is done with Python
- ✦ Scapy is used for creating the Ethernet packets and sending/receiving them in the network
- ✦ Registry databases are created with MySQL



Network Diagram

4 hosts, 3 PE (Provider Edge) nodes and 1 Provider

RE2EE Network Elements

- ✦ Host

- ✦ PE node

- ✦ Registry database

- ✦ REGISTRY and PROVIDERS tables

- ✦ Resolver

- ✦ Forwarder

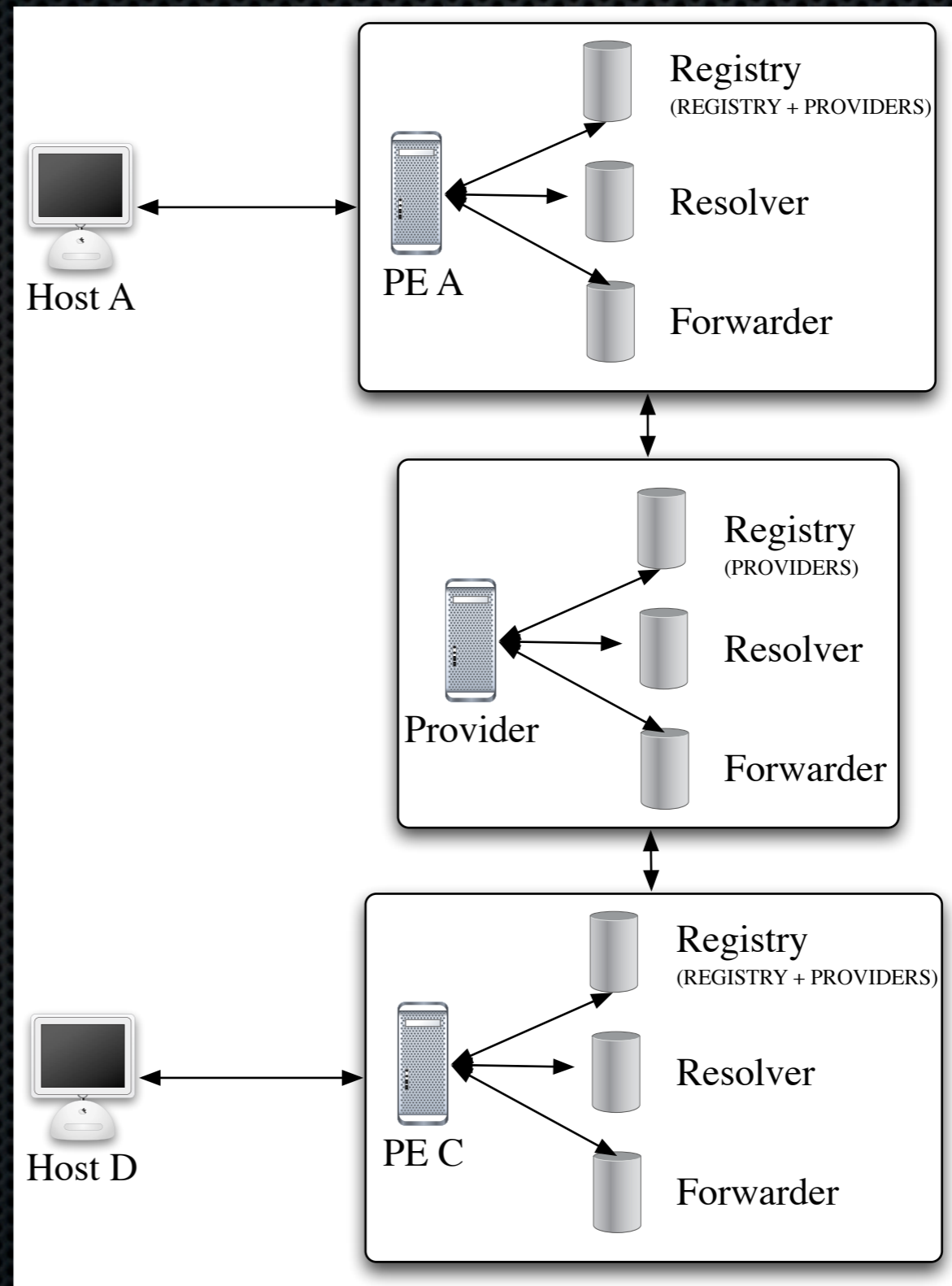
- ✦ Provider

- ✦ Registry database

- ✦ PROVIDERS table

- ✦ Resolver

- ✦ Forwarder



PoC Functionalities

- ✦ Home PE discovery by a host
- ✦ Host registering to Home PE
- ✦ Host sending data to other host
- ✦ Host receiving data from other host
- ✦ Inactivation of a host



Host A

```
find_PE("identity_of_service")  
dst_MAC="ff:ff:ff:ff:ff:ff"  
src_MAC="host_A_MAC"
```



PE A

```
response("result")  
dst_MAC="host_A_MAC"  
src_MAC="PE_A_MAC"
```

Home PE discovery by a host

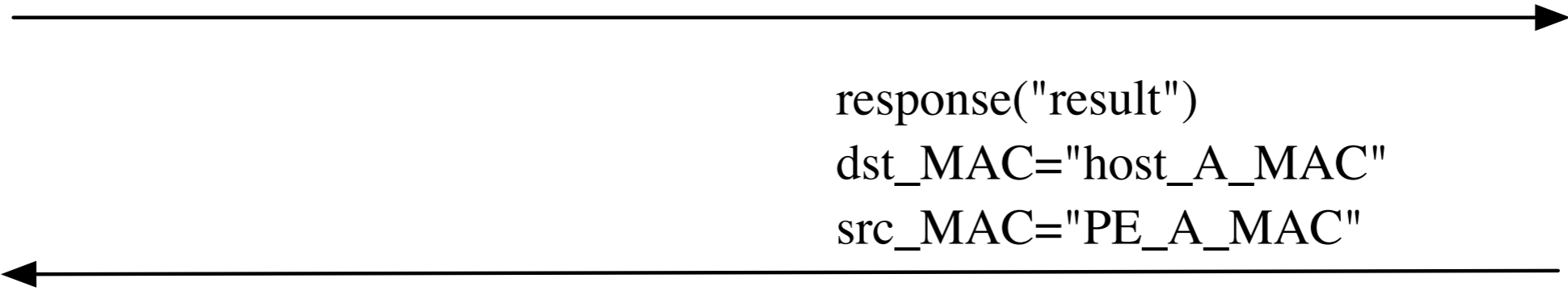


Host A

```
register("identity_of_host_A")  
dst_MAC="PE_A_MAC"  
src_MAC="host_A_MAC"
```



PE A



```
response("result")  
dst_MAC="host_A_MAC"  
src_MAC="PE_A_MAC"
```

Host registering to Home PE



Host A



PE A



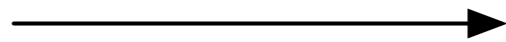
Host B

1.

```
send_packet("message","identity_of_host_B")
```

```
dst_MAC="PE_A_MAC"
```

```
src_MAC="host_A_MAC"
```



2.

```
forward("")
```

```
dst_MAC="host_B_MAC"
```

```
src_MAC="PE_A_MAC"
```

```
payload="identity_of_host_A"
```



Host A sending a message to Host B under the same home PE node



Host A



PE A



PE B



Host C

1.

`send_packet("message","identity_of_host_C")`

`dst_MAC="PE_A_MAC"`

`src_MAC="host_A_MAC"`



2.

`dst_MAC="PE_B_MAC"`

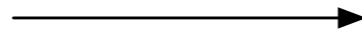
`src_MAC="PE_A_MAC"`

`encapsulate("")`

`dst_MAC="host_C_MAC"`

`src_MAC="PE_B_MAC"`

`payload="identity_of_host_A"`



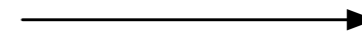
3.

`decapsulate and forward("")`

`dst_MAC="host_C_MAC"`

`src_MAC="PE_B_MAC"`

`payload="identity_of_host_A"`



Host A sending a message to Host C via PE A and PE B



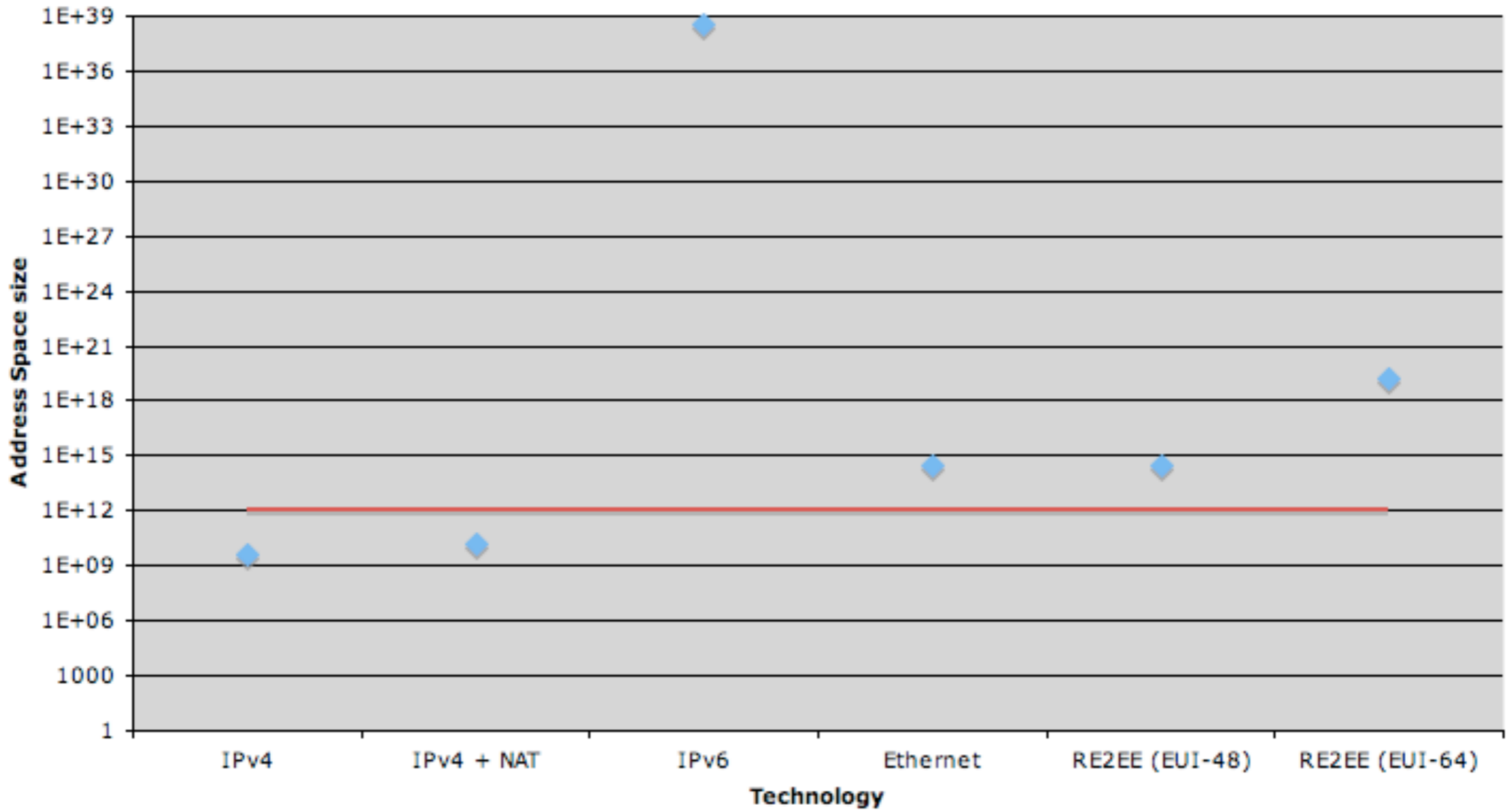
1.
send_packet("message","identity_of_host_D")
dst_MAC="PE_A_MAC"
src_MAC="host_A_MAC"
→
2.
dst_MAC="Provider_MAC"
src_MAC="PE_A_MAC"
encapsulate("")
dst_MAC="host_D_MAC"
src_MAC="PE_C_MAC"
payload="identity_of_host_A"
→
3.
forward("")
dst_MAC="PE_C_MAC"
src_MAC="Provider_MAC"
dst_MAC="host_D_MAC"
src_MAC="PE_C_MAC"
payload="identity_of_host_A"
→
4.
decapsulate and forward("")
dst_MAC="host_D_MAC"
src_MAC="PE_C_MAC"
payload="identity_of_host_A"
→

Host A sending a message to Host D via PE A, Provider and PE C

Results ^(1/2)

- ✦ There is a need for a new protocol in Internet for transmitting packets
- ✦ It needs to have the following features:
 - ✦ Large address space
 - ✦ Routing
 - ✦ Scalability
 - ✦ Security

Address Space sizes



Address Space sizes

Results (2/2)

✦ Routing

- ✦ Addresses need to have global hierarchy
- ✦ Ethernet has only Spanning Tree Protocol for routing -> need for better routing protocol like IS-IS

✦ Scalability

- ✦ One Routed Ethernet Provider Edge device can serve a minimum of 50 000 users
- ✦ There are about 1 billion Internet users all around the world
- ✦ Less than 20 000 devices are needed to serve 1 billion users

Results from the PoC

- ✦ PoC showed that it is possible to build a RE2EE network with small modifications
- ✦ No IP addresses were used
- ✦ Ethernet with registry database is enough for forwarding the traffic
- ✦ Using only identities for communicating hides the network from the hosts

Conclusions

- ✦ Routed End-to-End Ethernet would provide a long lasting solution to the problems, which we are having at the moment in the Internet
- ✦ RE2EE solution provides large enough address space, hierarchy, routing, security and scalability
- ✦ It is possible to use the same hosts in RE2EE network as in IP network

Future Research

- ✦ Make RE2EE network to work with IP networks
- ✦ Exact form of NSAP address
- ✦ Implement IS-IS routing
- ✦ Service discovery for different services
- ✦ Mobility management

Questions?